

All Variations on Perfectly Orderable Graphs

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An ordered graph is a graph whose vertices are positive integers. Two ordered graphs are isomorphic if the order-preserving bijection between their sets of vertices is a graph isomorphism. We identify the family of all sets S of ordered graphs with the following properties: (1) Each member of S is a P_4 (defined as a chordless path with four vertices and three edges). (2) If an ordered graph Z has no induced subgraph isomorphic (as an ordered graph) to a member of S , then Z is perfect. This work is related to Berge's Strong Perfect Graph Conjecture and was motivated by Chvátal's theorem on perfectly orderable graphs. © 1988 Academic Press, Inc.

1. THE RESULTS

Claude Berge defined a graph G to be *perfect* if for every induced subgraph H of G , the chromatic number $\chi(H)$ of H equals the largest number $\omega(H)$ of vertices in a clique of H .

Chvátal [1] introduced a class of perfect graphs, called *perfectly orderable graphs* and characterized by the existence of a linear order $<$ on the set of vertices such that no chordless path with vertices a, b, c, d and edges ab, bc, cd has $a < b$ and $d < c$.

This paper reports the results of a search for all variations on perfectly orderable graphs in a sense that we are about to make precise.

By an *ordered graph* we shall mean a graph whose vertices are distinct positive integers; two ordered graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ will be called *isomorphic* if there exists a bijection $f: V_1 \rightarrow V_2$ such that

$$f(i) < f(j) \quad \text{if and only if} \quad i < j$$

and such that

$$f(i)f(j) \in E_2 \quad \text{if and only if} \quad ij \in E_1.$$

By the *trace* of an ordered graph Z , we shall mean the subset S of the set of the 12 ordered graphs in Fig. 1 such that $X \in S$ if and only if Z contains

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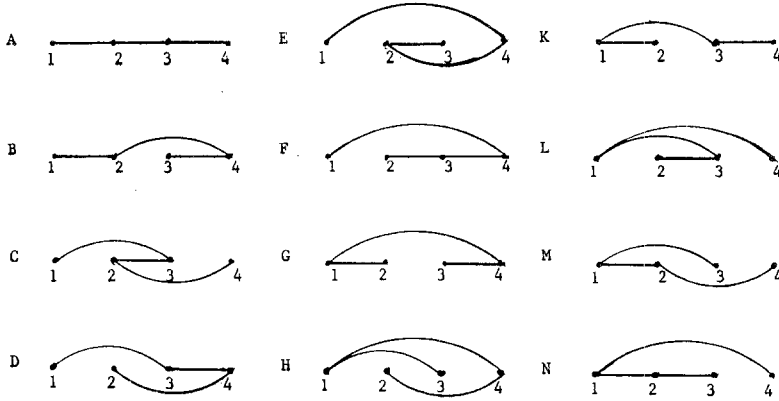


FIGURE 1

an induced ordered subgraph isomorphic to X . To simplify notation, we shall omit the braces when describing subsets of $\{A, B, \dots, N\}$: thus we shall write BDF for $\{B, D, F\}$ and so on.

Now Chvátal's theorem may be restated as follows:

if the trace of an ordered graph Z is disjoint from BDF then Z is perfect.

Our aim is to identify *all* sets S that could replace BDF in this theorem. The main results of this work involve the following sets:

$S_1 = ACEFK$	$S_2 = ACEMN$	$S_3 = CEGK$
$S_4 = ABHL$	$S_5 = ACEKN$	$S_6 = ABCEGH$
$S_7 = ACEKM$	$S_8 = AEGKM$	$S_9 = CDFG$
$S_{10} = ACDEK$	$S_{11} = ADEGK$	$S_{12} = ABCLN$
$S_{13} = ACDFL$	$S_{14} = BCGL$	$S_{15} = AEHK$
$S_{16} = ABCFL$	$S_{17} = ACGHKL$	$S_{18} = ABCDL$
$S_{19} = ABDGL$	$S_{20} = CGMN$	$S_{21} = ABCLM$
$S_{22} = ABGLM$	$S_{23} = EGHKM$	$S_{24} = DFGHK$
$S_{25} = DEG HK$	$S_{26} = EFGHK$	$S_{27} = CEFHK$
$S_{28} = EGHKN$	$S_{29} = CEHKN$	$S_{30} = BDGHL$
$S_{31} = BGHMN$	$S_{32} = BGHLM$	$S_{33} = BGHLN$
$S_{34} = BCHLN$	$S_{35} = BFGHL$	$S_{36} = BCFHL$
$S_{37} = BDF$	$S_{38} = KMN$	$S_{39} = LMN$
$S_{40} = DEF$	$S_{41} = ABCG HK$	$S_{42} = ACEGHL$

THEOREM 1. *If the trace of an ordered graph Z is disjoint from some S_i with $i = 1, 2, \dots, 42$ then Z is perfect.*

(To anticipate, the statements for $i = 1, 2, \dots, 36$ are a consequence of Hayward's theorem [3] that weakly triangulated graphs are perfect; the statements for $i = 37, 38, 39, 40$ are equivalent formulations of Chvátal's theorem [1]. Finally, the statements for $i = 41, 42$ follow from the perfection of opposition graphs which we prove in Theorem 2.)

Given a graph $G = (V, E)$ and a linear order $<$ on the set of its vertices, we shall write $u \rightarrow v$ if and only if $uv \in E$ and $u < v$. A P_4 (defined to be a chordless path with four vertices) involving vertices u, v, w, z and edges uv, vw, wz will be said to be *bad* if $u \rightarrow v$ and $w \rightarrow z$. As usual we shall let $N(u)$ stand for the set of neighbours of u in G ; \bar{G} will denote the complement of G .

A part of our argument used in the proof of Theorem 1 is of an independent interest; we shall set it on its own as Theorem 2. This theorem involves the notion of *opposition graph*: this is a graph with a linear order $<$ on the set of its vertices such that no induced P_4 is bad.

A *star-cutset* in a graph G is a nonempty set of vertices S such that some vertex $w \in S$ is adjacent to all remaining vertices in S . We shall say that the star-cutset is *centered* at w .

A graph G is said to be *minimal imperfect* if G is not perfect, but every proper induced subgraph H of G is perfect. We shall rely on the following result (see Chvátal [2]):

THE STAR-CUTSET LEMMA. *No minimal imperfect graph has a star-cutset.*

THEOREM 2. *An opposition graph G is either bipartite or else its complement \bar{G} has a star-cutset.*

Our two theorems will be proved in reverse order.

2. PROOFS

LEMMA 1. *If in an opposition graph $G = (V, E)$ there exist vertices a, c such that*

- (i) $ac \notin E, N(a) \cap N(c) \neq \emptyset$,
- (ii) $c \rightarrow x$ whenever $x \in N(a) \cap N(c)$,
- (iii) $a \rightarrow y$ whenever $y \in N(a) - N(c)$,

then \bar{G} has a star-cutset.

Proof of Lemma 1. Note first that if $N(a) - N(c) = \emptyset$ then trivially

$V - (\{c\} \cup N(c))$ is a star-cutset in \bar{G} centered at a . If $N(a) - N(c) \neq \emptyset$ then note that for every vertex u in $N(a) - N(c)$ and any vertex v in $N(a) \cap N(c)$ we must have $uv \in E$, for otherwise $cvau$ would be a bad P_4 . But now $V - N(a)$ is a star-cutset in \bar{G} centered at a . ■

A vertex s will be said to be a *source* if $w \rightarrow s$ for no $w \in V$; a vertex s will be said to be a *sink* if $s \rightarrow w$ for no $w \in V$; a vertex will be said to be *mixed* if it is neither a source nor a sink.

Proof of Theorem 2. Assume the theorem true for graphs with fewer vertices than G . If the graph is bipartite then there is nothing to prove, else there must exist a nonbipartite component G' of G . If this component has fewer vertices than G , we are done by the induction hypothesis: every star-cutset of \bar{G}' extends into a star-cutset of \bar{G} . Thus we may assume that G is connected and not bipartite.

Now the following algorithm finds a star-cutset in \bar{G} .

Let $S = \{a_1, a_2, \dots, a_k\}$ be the set of sources.

Step 1. If there are two sources a_i, a_j ($i \neq j$) and a vertex u such that $a_i \rightarrow u, a_j \rightarrow u$ then Lemma 1 applies with $a = a_i, c = a_j$, and so \bar{G} has a star-cutset.

Step 2. Write $w \in P_i$ iff $a_i \rightarrow w$ for $i = 1, 2, \dots, k$. Since we did not stop in Step 1, $P_i \cap P_j = \emptyset$ for all $i \neq j$. Write

$$P = \bigcup_{i=1}^k P_i.$$

If $uv \in E$ for some u in P_i, v in P_j with $i \neq j$ then stop: \bar{G} has a star-cutset. [To prove this claim, find the first vertex u in the linear order $<$ for which there are an edge $uv \in E$ and different subscripts i, j such that $u \in P_i, v \in P_j$. Now apply Lemma 1 with $c = u, a = a_j$.]

Step 3. Write $Q = V - (S \cup P)$. If there exists a vertex $w \in P_j$ such that $N(w) \cap Q = \emptyset$ then stop: \bar{G} has a star-cutset. [To see this, note that we must have $P_j - \{w\} \neq \emptyset$ for otherwise G would be disconnected, contradicting our assumption. Since G is connected, it follows that $N(w) \subseteq N(a_j)$ and $V - (\{a_j\} \cup (P_j - \{w\}))$ is a star-cutset in \bar{G} , centered at w .]

Step 4. If Q consists of sinks only, then let C be a component of P_i consisting of at least two vertices. (Such a component must exist or else G would be bipartite.)

4.1. If each $s \in Q$ is adjacent to either all the vertices of C or to none of them, then stop: for any $x \in C$ there is a star-cutset in \bar{G} consisting of x together with all vertices in $V - C$ that are not adjacent to x in G .

4.2. Now we may assume that some $s \in Q$ is adjacent to some, but not all, vertices in C . Since C is connected in G , it follows that there are adjacent vertices x, y in C with $x \rightarrow s, ys \notin E$. Since we did not stop in Step 3, there exists a vertex $t \in Q$ with $y \rightarrow t$. If $x \rightarrow t$ then stop: Lemma 1 with all the directions reversed and with $a = s, c = t$ guarantees that \bar{G} has a star-cutset.

Now $xt \notin E$. Note that we must have $N(s) \subseteq C$ or $N(t) \subseteq C$, for otherwise there would be vertices u, v in $P - C$ with $u \rightarrow s, v \rightarrow t$ and so $usxy$ or $vtpx$ would be a bad P_4 . But now $V - (\{a_i\} \cup P_i)$ is a star-cutset in \bar{G} , centered at r for $r = s$ or $r = t$.

Step 5. Now there exist mixed vertices in Q . Let z stand for the first mixed vertex in Q in the linear order $<$. Write $Z_1 = N(z) \cap Q$ and $Z_2 = N(z) \cap P$.

By minimality of z we have $Z_2 \neq \emptyset$.

5.1. If $Z_1 = \emptyset$ then $N(z) \subseteq N(a_j)$ for some j , for otherwise since z is a mixed vertex we find vertices $u \in P_i, v \in P_j$ such that $z \rightarrow u$ and $v \rightarrow z$, implying that a_jvzu is a bad P_4 , a contradiction. Clearly $V - (\{a_j\} \cup P_j)$ is a star-cutset in \bar{G} , centered at z .

5.2. Now $Z_1 \neq \emptyset$. We claim that every vertex $u \in Z_1$ is adjacent to every vertex $v \in Z_2$, for otherwise $uzva_j$, with $a_j \rightarrow v$, would be a bad P_4 . But now \bar{G} has a star-cutset consisting of $V - N(z)$ and centered at z . ■

LEMMA 2. *Every trace of a cycle of length at least five meets all of the following sets:*

$$\begin{array}{lll} W_1 = LMN & W_4 = CEFK & W_7 = ABHL \\ W_2 = GHK & W_5 = CMN & W_8 = ABLM \\ W_3 = ABDL & W_6 = CFGK & W_9 = CEKN. \end{array}$$

Proof of Lemma 2. Let T be the trace of a cycle of length at least five. We may assume that the vertices of the cycle are 1, 2, ..., n .

1. $T \cap W_1 \neq \emptyset$. Let i, j ($i < j$) be the two neighbours of 1 and let k be the neighbour of i distinct from 1. If $1 < k < i$ then $L \in T$; if $i < k < j$ then $N \in T$; if $j < k$ then $M \in T$.

2. $T \cap W_2 \neq \emptyset$. Let i, j ($i < j$) be the two neighbours of 1 and let k be the neighbour of j distinct from 1. If $1 < k < i$ then $H \in T$; if $i < k < j$ then $G \in T$; if $j < k$ then $K \in T$.

3. $T \cap W_3 \neq \emptyset$. Let i, j ($i < j$) be the two neighbours of 1 and let k be the neighbour of i distinct from 1. If $1 < k < i$ then $L \in T$. Thus we may

assume $i < k$. Now let t be the neighbour of k distinct from i . If $1 < t < i$ then $D \in T$; if $i < t < k$ then $B \in T$; if $k < t$ then $A \in T$.

4. $T \cap W_4 \neq \emptyset$. Let i, j ($i < j$) be the two neighbours of 1 and let k be the neighbour of j distinct from 1. If $j < k$ then $K \in T$. Thus we may assume $1 < k < j$. Now let t be the neighbour of k distinct from j . If $1 < t < k$ then $F \in T$; if $k < t < j$ then $E \in T$; if $j < t$ then $C \in T$.

5. $T \cap W_5 \neq \emptyset$. Let k be the smallest positive integer such that k is not a source; let i, j ($i < j$) be the two neighbours of k . If k is a sink then let t be the neighbour of j distinct from k ; by minimality of k , we have $k < t$ and so $C \in T$. If k is a mixed vertex then let t be the neighbour of i distinct from k . By minimality of k , we have $k < t$; if $k < t < j$ then $M \in T$; if $j < t$ then $N \in T$.

6. $T \cap W_6 \neq \emptyset$. Let k, i, j be as in the proof of $T \cap W_5 \neq \emptyset$; again, if k is a sink then $C \in T$. If k is mixed, then let x be the neighbour of i distinct from k . By minimality of k , we have $k < x$. Next, let y be the neighbour of x distinct from i . If $y < i$ then $E \in T$; if $k < y < x$ then $G \in T$; if $x < y$ then $K \in T$. Hence we may assume $i < y < k$. Finally, let z be the neighbour of y distinct from x . By minimality of k , we have $k < z$; if $k < z < x$ then $E \in T$; if $x < z$ then $C \in T$.

7. $T \cap W_7 \neq \emptyset$. Let k, i, j be as in the proof of $T \cap W_5 \neq \emptyset$. If k is a sink then let t be the neighbour of i distinct from k ; by minimality of k , we have $k < t$, and so $L \in T$. If k is mixed then let x be the neighbour of j distinct from k . If $k < x < j$ then $B \in T$; if $j < x$ then $A \in T$. Hence we may assume $x < k$. Let y be the neighbour of x distinct from j . By minimality of k , we have $k < y$. If $k < y < j$ then $H \in T$; if $j < y$ then $L \in T$.

8. $T \cap W_8 \neq \emptyset$. Let k be the smallest positive integer such that k is a sink; let i, j ($i < j$) be the two neighbours of k and let x be a neighbour of i distinct from k . If $x < i$ then $B \in T$; if $k < x$ then $L \in T$. Hence we may assume $i < x < k$. Let y be the neighbour of x distinct from i . By minimality of k , we have $x < y$; if $k < y$ then $M \in T$. Hence we may assume $x < y < k$. Finally, let z be the neighbour of y distinct from x . By minimality of k , we have $y < z$, and so $A \in T$.

9. $T \cap W_9 \neq \emptyset$. Let k, i, j be as in the proof of $T \cap W_8 \neq \emptyset$, and let x be the neighbour of j distinct from k . If $j < x < k$ then $E \in T$; if $k < x$ then $C \in T$. Hence we may assume $x < j$.

Set $v_0 = k$, $v_1 = j$, $v_2 = x$, $t = 2$ and, as long as v_t has a neighbour y with $y < v_t$, keep setting $v_{t+1} = y$, $t \leftarrow t + 1$. When this process terminates v_t has a neighbour y with $y \neq v_{t-1}$, $v_t < y$. If $v_t < y < v_{t-1}$ then $K \in T$; if $v_{t-2} < y$ then $N \in T$. Hence we may assume that $v_{t-1} < y < v_{t-2}$; note that $v_{t-2} < k$.

Finally, let z be the neighbour of y distinct from v_i . By minimality of k , we have $y < z$, and so $K \in T$. ■

For each ordered graph X with vertices $1, 2, \dots, n$ we shall denote by

$F(X)$ the ordered graph resulting from X when each vertex i is relabeled as $n+1-i$.

$C(X)$ the ordered graph resulting from X when every two non-adjacent vertices are made adjacent and vice versa.

On the set of the 12 ordered graphs in Fig. 1 these two mappings act as shown in Fig. 2.

Note that, for every unordered graph Q and for each subset S of $\{A, B, \dots, N\}$, the following three statements are equivalent:

(i) S meets the trace of every ordered graph isomorphic (as an unordered graph) to Q ,

(ii) $F(S)$ meets the trace of every ordered graph isomorphic (as an unordered graph) to Q ,

(iii) $C(S)$ meets the trace of every ordered graph isomorphic (as an unordered graph) to the complement \bar{Q} of Q .

For further reference, observe that the family S_1, S_2, \dots, S_{42} is "closed" under applications of the mappings C and F : more precisely, for each S_i there are S_j, S_k such that $F(S_i) = S_j$ and $C(S_i) = S_k$ with $i, j, k = 1, 2, \dots, 42$ (see Fig. 3).

A graph G is called *weakly triangulated* if it contains no induced subgraph isomorphic to a C_k or \bar{C}_k with $k \geq 5$. Hayward [3] has proved that weakly triangulated graphs are perfect.

LEMMA 3. *Each of the sets S_i with $i = 1, 2, \dots, 36$ has the following property: if the trace of an ordered graph Z is disjoint from S_i then Z is weakly triangulated.*

Proof of Lemma 3. Clearly, we only need prove that Z contains no induced subgraph isomorphic to a C_k or \bar{C}_k with $k \geq 5$.

X	A	B	C	D	E	F	G	H	K	L	M	N
$F(X)$	A	K	C	M	L	N	G	H	B	E	D	F
$C(X)$	H	L	G	N	K	M	C	A	E	B	F	D

FIGURE 2

S_i	$F(S_i)$	$C(S_i)$	S_i	$F(S_i)$	$C(S_i)$
S_1	S_{12}	S_{23}	S_{22}	S_{11}	S_{36}
S_2	S_{13}	S_{24}	S_{23}	S_{30}	S_1
S_3	S_{14}	S_3	S_{24}	S_{31}	S_2
S_4	S_{15}	S_4	S_{25}	S_{32}	S_5
S_5	S_{16}	S_{25}	S_{26}	S_{33}	S_7
S_6	S_{17}	S_{17}	S_{27}	S_{34}	S_8
S_7	S_{18}	S_{26}	S_{28}	S_{35}	S_{10}
S_8	S_{19}	S_{27}	S_{29}	S_{36}	S_{11}
S_9	S_{20}	S_9	S_{30}	S_{23}	S_{12}
S_{10}	S_{21}	S_{28}	S_{31}	S_{24}	S_{13}
S_{11}	S_{22}	S_{29}	S_{32}	S_{25}	S_{16}
S_{12}	S_1	S_{30}	S_{33}	S_{26}	S_{18}
S_{13}	S_2	S_{31}	S_{34}	S_{27}	S_{19}
S_{14}	S_3	S_{14}	S_{35}	S_{28}	S_{21}
S_{15}	S_4	S_{15}	S_{36}	S_{29}	S_{22}
S_{16}	S_5	S_{32}	S_{37}	S_{38}	S_{39}
S_{17}	S_6	S_6	S_{38}	S_{37}	S_{40}
S_{18}	S_7	S_{33}	S_{39}	S_{40}	S_{37}
S_{19}	S_8	S_{34}	S_{40}	S_{39}	S_{38}
S_{20}	S_9	S_9	S_{41}	S_{41}	S_{42}
S_{21}	S_{10}	S_{35}	S_{42}	S_{42}	S_{41}

FIGURE 3

The cases of $i = 1, 2, \dots, 11$ may be settled by observing that

$$\begin{aligned}
 S_1 &= W_4 \cup C(W_2) & S_7 &= F(W_3) \cup C(W_2) \\
 S_2 &= W_5 \cup C(W_2) & S_8 &= F(W_3) \cup C(W_4) \\
 S_3 &= W_6 \cup C(W_6) & S_9 &= F(W_5) \cup C(W_5) \\
 S_4 &= W_7 \cup C(W_7) & S_{10} &= F(W_8) \cup C(W_2) \\
 S_5 &= W_9 \cup C(W_2) & S_{11} &= F(W_8) \cup C(W_9). \\
 S_6 &= F(W_2) \cup C(W_2)
 \end{aligned}$$

The cases of $i = 12, 13, \dots, 22$ may be settled by observing that

$$\begin{aligned}
 S_{12} &= F(S_1) & S_{18} &= F(S_7) \\
 S_{13} &= F(S_2) & S_{19} &= F(S_8) \\
 S_{14} &= F(S_3) & S_{20} &= F(S_9) \\
 S_{15} &= F(S_4) & S_{21} &= F(S_{10}) \\
 S_{16} &= F(S_5) & S_{22} &= F(S_{11}). \\
 S_{17} &= F(S_6)
 \end{aligned}$$

Finally, the cases of $i = 23, 24, \dots, 36$ may be settled by observing that

$$\begin{array}{ll} S_{23} = C(S_1) & S_{30} = C(S_{12}) \\ S_{24} = C(S_2) & S_{31} = C(S_{13}) \\ S_{25} = C(S_5) & S_{32} = C(S_{16}) \\ S_{26} = C(S_7) & S_{33} = C(S_{18}) \\ S_{27} = C(S_8) & S_{34} = C(S_{19}) \\ S_{28} = C(S_{10}) & S_{35} = C(S_{21}) \\ S_{29} = C(S_{11}) & S_{36} = C(S_{22}). \end{array}$$

Proof of Theorem 1. Note that for sets S_i with $i = 1, 2, \dots, 36$ the desired conclusion follows immediately from Lemma 3 and Hayward's result.

Next, for sets S_i with $i = 37, 38, 39, 40$ the conclusion follows as a corollary to Chvátal's theorem.

Finally, to complete the proof we shall rely on the following corollary of Theorem 2:

Opposition graphs are perfect.

[Assuming the contrary, a minimal counterexample would be minimal imperfect. By Theorem 2 and the star-cutset lemma it follows that \bar{G} is perfect. Now by a fundamental result of Lovász [4], G is perfect if and only if \bar{G} is perfect, a contradiction.]

Now we only need observe that if the trace of an ordered graph Z is disjoint from S_{42} , then Z is an opposition graph, and that $C(S_{42}) = S_{41}$. This settles the cases of S_i with $i = 41, 42$. ■

We shall say that a graph G is *breakable* (a term coined by V. Chvátal) if at least one of G, \bar{G} has a star-cutset.

Chvátal [2] proved that if a graph G is perfectly orderable, then it is either bipartite or else \bar{G} has a star-cutset.

Hayward [3] proved that if a graph is weakly triangulated then it is breakable.

These facts combined with Theorem 2 yield the following proposition:

If the trace of an ordered graph Z is disjoint from some S_i with $i = 1, 2, \dots, 42$ then Z is breakable or else Z or \bar{Z} is bipartite.

By virtue of the Star-Cutset Lemma, this proposition implies Theorem 1.

We note that the converse of Theorem 1 is false: there exist graphs which are perfect but have no ordering which is one of these variations on perfectly orderable graphs. One such graph H is featured in Fig. 4.

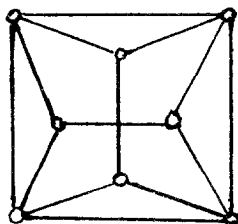


FIGURE 4

To justify this claim it suffices to note that the graph H is perfect, neither H nor \bar{H} is bipartite and that H is not breakable.

To prove the completeness of the list S_1, S_2, \dots, S_{42} we prove that every set that could play the role of some S_i with $i = 1, 2, \dots, 42$ in Theorem 1 must contain some S_j with $j = 1, 2, \dots, 42$. The proof itself is long and tedious but the basic idea is quite simple: we identify a family of sets $T_j, j = 1, 2, \dots, 118$ such that each T_j is either the trace of an odd cycle or the trace of the complement of an odd cycle. We then prove that if a subset S of $AB \dots N$ meets each T_j with $j = 1, 2, \dots, 118$, then S contains some S_i with $i = 1, 2, \dots, 42$.

The interested reader can find the details in [5].

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